

科目：計算理論

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1. (30%) Let $\Sigma = \{0, 1\}$, $w \in \Sigma^*$, w^R denote the reverse of w , and $\#_a(w)$ denote the number of symbol a in string w , where $a \in \Sigma$. Please prove or disprove the following statements:

- (a) (5%) $L_{1a} = \{w \mid \#_0(w) \text{ is even and } \#_1(w) \text{ is odd}\}$ is regular.
 (b) (5%) $L_{1b} = \{w \mid \#_0(w) = \#_1(w)\}$ is regular.
 (c) (10%) $L_{1c} = \{w \mid \#_0(w) = \#_1(w)\}$ is context free.
 (d) (10%) $L_{1d} = \{w \mid w = w^R \text{ and } \#_0(w) = \#_1(w)\}$ is context free.

2. (10%) If A is a set of natural numbers and k is a natural number greater than 1, let

$$B_k(A) = \{w \mid w \text{ is the representation in base } k \text{ of some number in } A\}.$$

Here, we do not allow leading 0s in the representation of a number. For example, $B_2(\{3, 5\}) = \{11, 101\}$ and $B_3(\{3, 5\}) = \{10, 12\}$. Give an example of a set A for which $B_2(A)$ is regular but $B_3(A)$ is not regular. Prove that your example works.

3. (10%) If A and B are languages, define

$$A \diamond B = \{xy \mid x \in A \text{ and } y \in B \text{ and } |x| = |y|\}.$$

part A Show that if A and B are regular, then $A \diamond B$ is context free.

- part B 4. (10%) Prove that a language is decidable if and only if it is Turing-recognizable (r.e.) and co-Turing-recognizable (co-r.e.).

5. (10%) Language A is *mapping reducible* to language B , written $A \leq_m B$, if there is a computable function $f : \Sigma^* \rightarrow \Sigma^*$, where for every w , $w \in A \iff f(w) \in B$. Prove that if $A \leq_m B$ and B is decidable, then A is decidable.

6. (10%) Define $L = \{\langle M, N \rangle \mid M, N \text{ are Turing machines and } L(M) = L(N)\}$. Prove that L is neither Turing-recognizable nor co-Turing-recognizable.

7. (10%) A *2cnf-formula* is an AND of clauses, where each clause is an OR of at most two literals. Let $2SAT = \{\langle \phi \rangle \mid \phi \text{ is a satisfiable 2cnf-formula}\}$. Show that $2SAT \in P$.

8. (10%) Consider the following problem: Let S be a set of integers. Can we partition S into two parts: S_1 and S_2 such that the sum of integers in S_1 is equal to the sum of integers in S_2 ? Prove that this problem is NP-complete. You can use the problem *SUBSET-SUM* for reduction.