

科目：演算法(A)

日期：101年7月26日 第1頁共3頁

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\* 請將答案依題號順序寫入答案卷

答題時字跡需工整，否則不予計分。Write your answers legibly; otherwise you will get zero score.

1. 16% Interval graphs, activity-selection problem.

Suppose we have a set  $S = \{a_1, a_2, \dots, a_n\}$  of  $n$  proposed activities. Each activity  $a_i$  has a start time  $s_i$  and a finish time  $f_i$ , where  $0 \leq s_i < f_i < \infty$ . If selected, activity  $a_i$  takes place during the half-open time interval  $[s_i, f_i)$ . Activities  $a_i$  and  $a_j$  are compatible if the intervals  $[s_i, f_i)$  and  $[s_j, f_j)$  do not overlap. That is,  $a_i$  and  $a_j$  are compatible if  $s_i \geq f_j$  or  $s_j \geq f_i$ . In the activity-selection problem, we wish to select a maximum-size subset of mutually compatible activities.

For the following greedy idea, pick out the ones that yield an optimal solution. You must give proof.

- (1) Selecting the first activity to start (that is compatible with all previously selected activities).
  - (2) Selecting the first activity to finish (that is compatible with all previously selected activities).
  - (3) Selecting the last activity to start (that is compatible with all previously selected activities).
  - (4) Selecting the last activity to finish (that is compatible with all previously selected activities).
  - (5) Selecting the activity of least duration (that is compatible with all previously selected activities).
  - (6) Selecting the compatible activity that overlaps the fewest other remaining activities.
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## 2. 17% Minimum spanning trees.

In this problem, we give pseudocode for three different algorithms. Each one takes a graph as input and returns a set of edges  $T$ . For each algorithm, you must either prove that  $T$  is a minimum spanning tree or prove that  $T$  is not a minimum spanning tree.

a、MAYBE-MST-A( $G, w$ )

```
1  sort the edge into nonincreasing order of edge weights  $w$ 
2   $T \leftarrow E$ 
3  for each edge  $e$ , taken in nonincreasing order by weight
4      do if  $T - \{e\}$  is a connected graph
5          then  $T \leftarrow T - e$ 
6  return  $T$ 
```

b、MAYBE-MST-B( $G, w$ )

```
1   $T \leftarrow \emptyset$ 
2  for each edge  $e$ , taken in arbitrary order
3      do if  $T \cup \{e\}$ 
4          then  $T \leftarrow T \cup e$ 
5  return  $T$ 
```

c、MAYBE-MST-C( $G, w$ )

```
1   $T \leftarrow \emptyset$ 
2  for each edge  $e$ , taken in arbitrary order
3      do  $T \leftarrow T \cup e$ 
4          if  $T$  has a cycle  $C$ 
5              then let  $e'$  be the maximum-weight edge on  $C$ 
6                   $T \leftarrow T - \{e'\}$ 
7  return  $T$ 
```

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3.

(A) 4% Let  $G$  be a directed graph with any kind of edge weights, positive and/or negative weights, having positive weight cycles and/or negative weight cycles. Assumes  $NP \neq P$ .

For each of the following problems classify whether it is NP-hard or P.

(1) Find a negative weight cycle.

(2) Find a negative weight cycle with the least absolute value total weight.

(3) Find a positive weight cycle.

(4) Find a positive weight cycle with the least total weight.

(B) 13% Prove your answers. (You may explain briefly, give a simple illustration, or give the name of the algorithm.)

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科目：演算法(B)

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答題時字跡需工整，否則不予計分。Write your answers legibly; otherwise you will get zero score.

- (15%) Rank the following functions by order of growth; that is, find an arrangement  $g_1, g_2, \dots, g_n$  of functions satisfying  $g_1 = \Omega(g_2)$ ,  $g_2 = \Omega(g_3)$ , ...,  $g_{n-1} = \Omega(g_n)$ . Partition your list into equivalence classes such that  $f(n)$  and  $g(n)$  are in the same class if and only if  $f(n) = \Theta(g(n))$ . Hint: derive the simplest formula for each function.
 

(a) $(\log n)^{\log 2n}$	(i) $(n!)!$
(b) $n (\log n) / (\log \log n)^3$	(j) $(n!)^n$
(c) $n^{\log (\log n)^2}$	(k) $1 \times (n-1) + 2 \times (n-2) + \dots + (n-1) \times 1$
(d) $\log((n \log n)^{n \log n})$	(l) $n + n/2^2 + n/3^2 \dots + n/(n^2)$
(e) $n^n$	(m) $2 \times (n^2-2) + 4 \times (n^2-4) + \dots + (n^2-2) \times 2$
(f) $(1) + (1 + 1/2) + (1 + 1/2 + 1/3)$ $+ \dots + (1 + 1/2 + \dots + 1/n)$	(n) $1 + 1/2 + 1/3 \dots + 1/(10^n)$
(g) $2^{3n}$	(o) $2^{\lg n}$
(h) $3^{2n}$	(p) $4^{\lg n}$
- (10%) In a hash table in which collisions are resolved by chaining, let the load fraction be  $\alpha = n/m$ . Prove that an unsuccessful search takes expected time  $\Theta(1 + \alpha)$ , under the assumption of simple uniform hashing.
- (13%) Given  $n$  integers  $(a_1, a_2, \dots, a_n)$ , where the number of distinct integers is  $O(\log n)$ , design a comparison sort algorithm using at most  $O(n \log \log n)$  comparisons. But, our textbook seems to say “the lower bound for sorting is  $\Omega(n \log n)$ ” (this seems contradictory to your algorithm). Explain why.
- (12%) Briefly describe the dynamic programming algorithm for optimal parenthesization of a matrix-chain product with the smallest number of scalar multiplications. Illustrate your algorithm for the matrix-chain product whose sequence of dimensions is  $\{6, 7, 3, 1, 2, 4, 5\}$ .